

## MIPS registers

MIPS has 32 registers r0-r31. The conventional use of these registers is as follows:

register	assembly name	Comment
r0	\$zero	Always 0
r1	\$at	Reserved for assembler
r2-r3	\$v0-\$v1	Stores results
r4-r7	\$a0-\$a3	Stores arguments
r8-r15	\$t0-\$t7	Temporaries, not saved
r16-r23	\$s0-\$s7	Contents saved for use later
r24-r25	\$t8-\$t9	More temporaries, not saved
r26-r27	\$k0-\$k1	Reserved by operating system
r28	\$gp	Global pointer
r29	\$sp	Stack pointer
r30	\$fp	Frame pointer
r31	\$ra	Return address

## Loading a 32-bit constant into a register

```
lui $s0, 42      # load upper-half immediate  
ori $s0, $s0, 18 # (one can also use andi)
```

What is the end result?

## Review the logical operations

Shift left logical      sll

Shift right logical     srl

Bit-by-bit AND          and, andi (and immediate)

```
sll $t2, $s0, 4      # register $t2 := register $s0 << 4
```

s0 = 0000 0000 0000 0000 0000 0000 0000 1001

t2 = 0000 0000 0000 0000 0000 0000 1001 0000



Op = 0	rs = 0	rt = 16	rd = 10	shamt = 4	function = 0
--------	--------	---------	---------	-----------	--------------

(s0 = r16, t2 = r10)

What are the uses of shift instructions?

Multiply or divide by some power of 2.

Implement general multiplication using addition and shift

## Making decisions

```
if (i == j)      f = g + h;   else    f = g - h
```

Use **bne** = branch-not-equal, **beq** = branch-equal, and **j** = jump

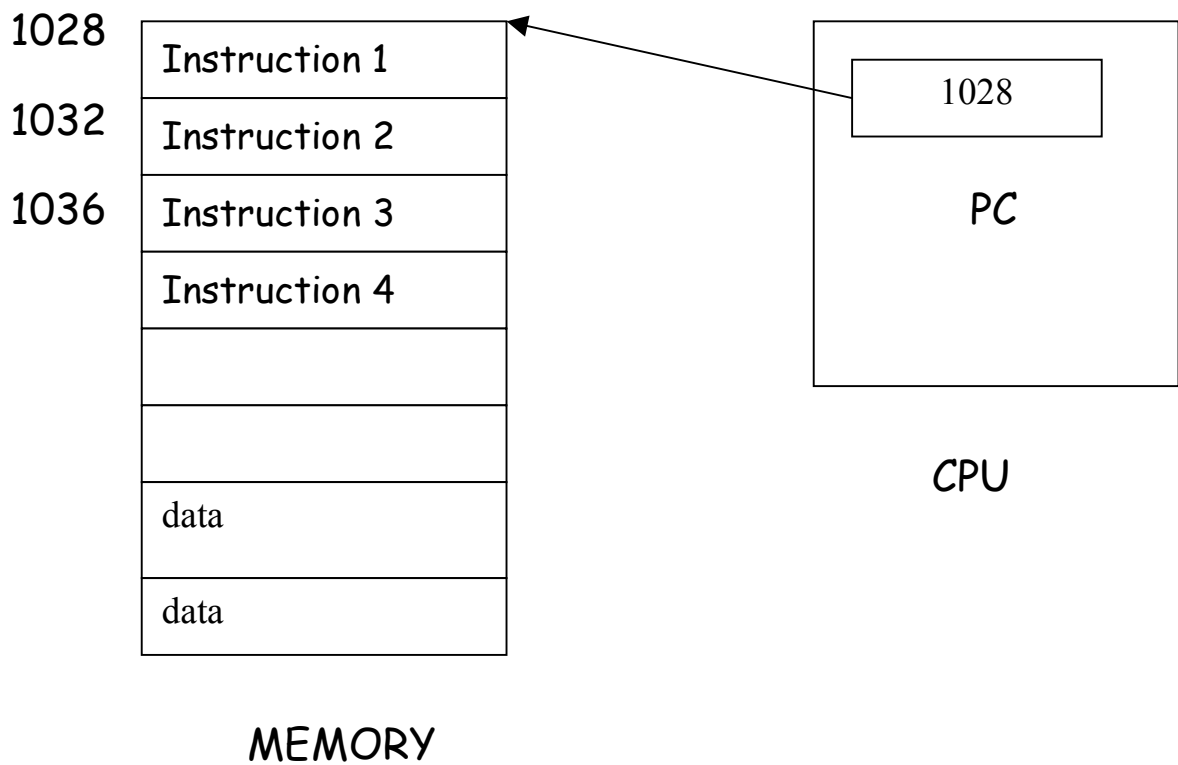
Assume that  $f, g, h$ , are mapped into  $\$s0, \$s1, \$s2$

$i, j$  are mapped into  $\$s3, \$s4$

```
        bne $s3, $s4, Else      # goto Else when i=j
        add $s0, $s1, $s2      # f = g + h
        j   Exit               # goto Exit
Else:   sub $s0, $s1, $s2      # f = g - h
Exit:
```

## The program counter

Every machine has a **program counter** (called PC) that points to the next instruction to be executed.



Ordinarily, PC is incremented by 4 after each instruction is executed. A branch instruction alters the flow of control by modifying the PC.

## Compiling a while loop

```
while (A[i] == k)    i = i + j;
```

Initially \$s3, \$s4, \$s5 contains i, j, k respectively.

Let \$s6 store the base of the array A. Each element of A is a 32-bit word.

Loop:	add \$t1, \$s3, \$s3	# \$t1 = 2*i
	add \$t1, \$t1, \$t1	# \$t1 = 4*i
	add \$t1, \$t1, \$s6	# \$t1 contains address of A[i]
	lw \$t0, 0(\$t1)	# \$t0 contains \$A[i]
	add \$s3, \$s3, \$s4	# i = i + j
	bne \$t0, \$s5, Exit	# goto Exit if A[i] ≠ k
	j Loop	# goto Loop
Exit:	<next instruction>	

Note the use of pointers.

## Exercise

Add the elements of an array  $A[0..63]$ . Assume that the first element of the array is stored from address 200. Store the sum in address 800.

## System Call

The program takes the help of the operating system to do some input or output operation. Example

```
li    $v0, 5          # System call code for Read Integer
syscall                # Read N into $v0
```

## Compiling a switch statement

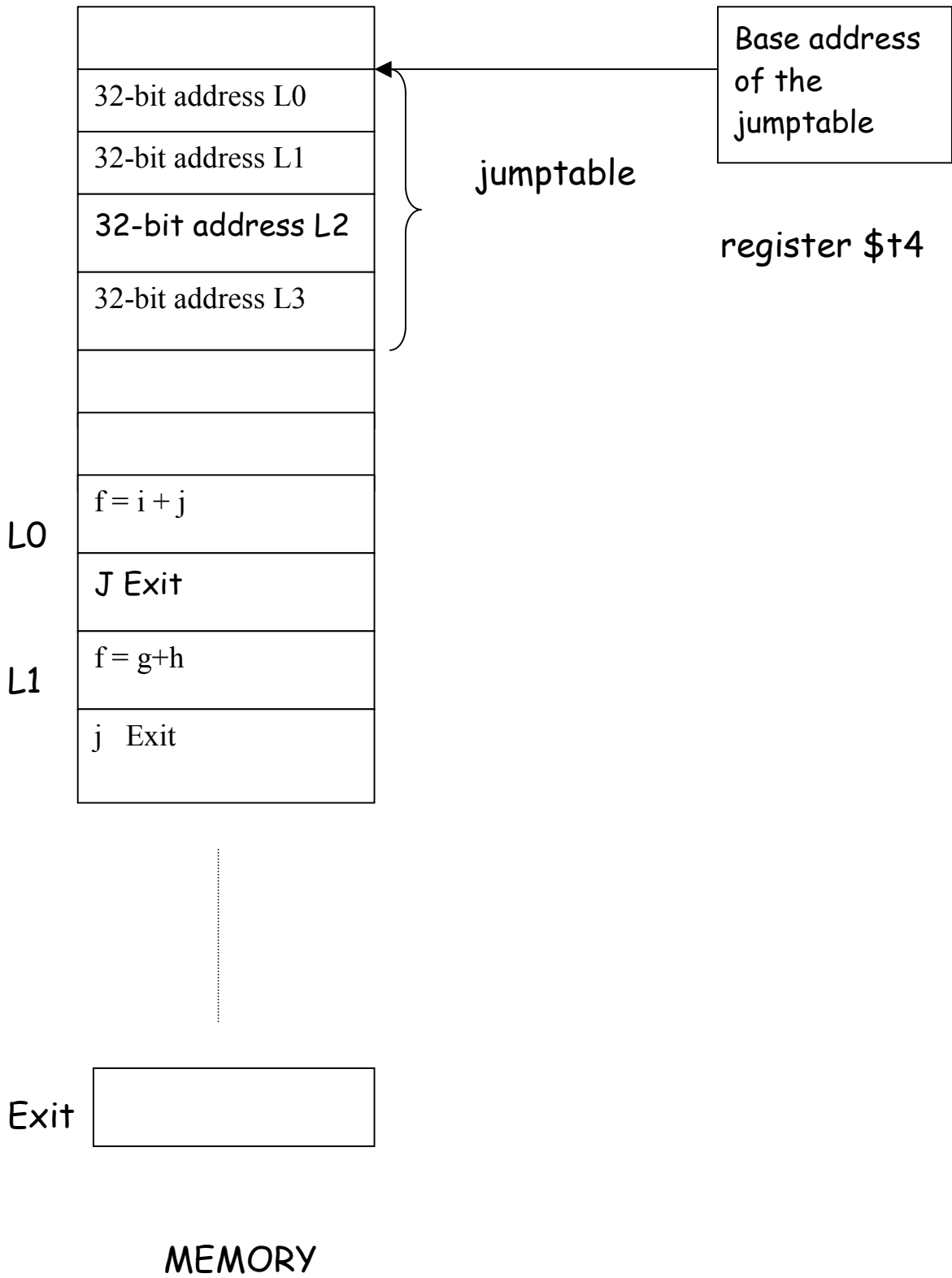
```
switch (k) {  
    case 0:  f = i + j; break;  
    case 1:  f = g + h; break;  
    case 2:  f = g - h; break;  
    case 3:  f = i - j; break;  
}
```

Assume, \$s0-\$s5 contain f, g, h, i, j, k.

Assume \$t2 contains 4.

```
slt $t3, $s5, $zero    # if k<0 then $t3 = 1 else $t3=0  
bne $t3, $zero, Exit  # if k<0 then Exit  
slt $t3, $s5, $t2     # if k<4 then $t3 = 1 else $t3=0  
beq $t3, $zero, Exit  # if k≥ 4 the Exit
```

What next? Jump to the right case!





Here is the remainder of the program;

```
add $t1, $s5, $s5      # t1 = 2*k
add $t1, $t1, $t1      # t1 = 4*k
add $t1, $t1, $t4      # t1 = base address + 4*k
lw $t0, 0($t1)         # load the address pointed
                        # by t1 into register t0
jr $t0                 # jump to addr pointed by t0
L0: add $s0, $s3, $s4   # f = i + j
    J Exit
L1: add $s0, $s1, $s2   # f = g+h
    J Exit
L2: sub $s0, $s1, $s2   # f = g-h
    J Exit
L3: sub $s0, $s3, $s4   # f = i - j
Exit: <next instruction>
```