

# 22C:153 Homework 1

Due: Tuesday, 2/25

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## On basic probability theory:

- (1) Problem C.2-9 (page 1106). Explain your answer by calculating the relevant probabilities.
- (2) Problem C.3-6 (page 1111). Just show this assuming that  $X$  is a discrete random variable that takes on finitely many non-negative values.

## On probabilistic analysis:

- (3) Problem 5.2-5 (page 99).
- (4) Problem 5.3-4 (page 105).

## On randomized quicksort:

- (5) Problem 7-2 (page 160).
- (6) Problem 7-4 (page 162).

## On randomized selection:

- (7) In the analysis of randomized selection we derived the following sum:

$$\sum_{i=1}^{k-2} \sum_{j=i+1}^{k-1} \frac{2}{(k-i+1)} + \sum_{i=k+1}^{n-1} \sum_{j=i+1}^n \frac{2}{(j-k+1)} + \sum_{i=1}^k \sum_{j=k}^n \frac{2}{(j-i+1)}.$$

Show that for any integer  $k$ ,  $1 \leq k \leq n$ , the above sum is  $O(n)$ .

## On Karger's min-cut algorithm:

- (8) Consider the modification of Karger's min-cut algorithm mentioned in class in which, instead of picking the endpoints of an edge to contract, we pick (uniformly at random) a pair of vertices and contract them. Show an example of a graph with  $n$  vertices such that the probability that the modified algorithm will find a min-cut in the graph is exponentially small (that is, the probability is bounded above by  $\Theta(1/e(n))$  where  $e(n)$  is an exponential function in  $n$ ).
- (9) We showed that the probability that Karger's min-cut algorithm does not find a min-cut is less than  $1/e$ . How many times does this algorithm have to be repeated if we want this probability to be less than  $1/100$ ? Justify your answer.
- (10) Describe a data structure to store a multigraph (a graph with multiple edges between pairs of vertices) such that each **CONTRACT** can be performed in  $O(n)$  time on an  $n$ -vertex graph. Furthermore, suppose that we start with the graph  $G = (V, E)$  and have contracted the endpoints of edges  $e_1, e_2, \dots, e_k$ . Let the resulting graph be called  $H$ . Each edge in  $H$  corresponds to an edge in the original graph  $G$ . Given an edge in  $H$  how quickly can you find the corresponding edge in  $G$  using your data structure?  
(**Hint:** Consider a small modification of an adjacency matrix.)