Solving Triangular Systems in Parallel

Summary of Forward Substitution Algorithms

- Sequential Inner Product Version
- Sequential Vector Sum
- Fine-Grain Parallel Pseudocode

if $i = j$
  recv sum reduction σ_i
  $x_i = (b_i - \sigma_i)/l_{ii}$
  broadcast x_i to tasks (k, i)
  $(k = i + 1, \cdots, n)$
else
  recv x_j
  $t = l_{i,j}x_j$
  send t for sum reduction across tasks
  (i, k) $(k = 1, \cdots, (i - 1))$ to task (i, i)
end

- Row Partitioning for Vector Sum Algorithm (Fan-out)

  The output loop (for all the columns) is sequential. Consequently the number of sequential steps is $O(n)$. Processors have to wait for information above them before they start updating b_j. Processors *broadcast* the x_j's they are responsible for, as soon as they are calculated.

  for $j = 1$ **to** n
    if $j \in myrows$ **then**
      $x_j = b_j / l_{jj}$
      broadcast x_j to other tasks
    else
      recv x_j
    end
    for $i \in myrows, \;\; i > j$
      $b_i = b_i - l_{ij} x_j$
    end
  end

- Column Partitioning for Inner Product Algorithm (Fan-in)

  The output loop (for all the rows) is sequential. Consequently the number of sequential steps is $O(n)$. For each row a partial inner product (t) is calculated and used in a *reduce* accross other processors. Processors have to wait for informationl on their left before calculating x_j's.

  for $i = 1$ **to** n
    $t = 0$
    for $j \in mycols, \;\; j < i$
      $t = t + l_{ij} x_j$
    end
    if $i \in mycols$ **then**
      recv sum reduction of t
      $x_i = (b_i - t)/l_{ii}$
    else
      send t for sum reduction across tasks
    end
  end

- Wavefront Vector Sum Algorithm

for $j \in mycols$
  for $k = 1$ **to** (# of $segments$)
    recv $segment$
    if $k = 1$ **then**
      $x_j = (b_j - z_j)/l_{jj}$
      $segment = segment - \{z_j\}$
    end
    for $z_i \in segment$
      $z_i = z_i + l_{ij}x_j$
    end
    if $|segment| > 0$ **then**
      send $segment$ to task with column $j + 1$
    end
  end
end

- Wavefront Scalar Product Algorithm

for $j \in myrows$
  for $k = 1$ **to** (# of $segments - 1$)
    recv $segment$
    send $segment$ to task owing row $i + 1$
    for $x_j \in segment$
      $b_i = b_i - l_{ij}x_j$
    end
  end
  recv $segment$
  for $x_j \in segment$
    $b_i = b_i - l_{ij}x_j$
  end
  $x_i = b_i/l_{ii}$
  $segment = segment \cup \{x_i\}$
  send $segment$ to task owing row $i + 1$
end