

Example program proof — while rule

In this example, we prove a program assertion establishing partial correctness of a program fragment computing the factorial of an integer.

```

    { N ≥ 0 }
M:= 0; F = 1;
while M < N do
begin M := M+1; F := M*F end
    { F = N! }

```

To construct the proof of this program, we need to determine an intermediate formula \mathbb{P} that will serve as a loop invariant. This will describe the relationship of the variable values at the intermediate point noted below.

```

    { N ≥ 0 }
M:= 0; F = 1;
     $\mathbb{P}$ 
while M < N do
begin M := M+1; F := M*F end
    { F = N! }

```

We take \mathbb{P} to be the formula $\{ F = M! \wedge 0 \leq M \leq N \}$. Then the proof is as follows:

1. By two applications of the axiom of assignment and the sequential execution rule
 - | $\{ 1 = 0! \wedge 0 \leq N \} M:= 0; F = 1 \mathbb{P}$
 and the pre-condition is logically equivalent to the pre-condition of the program.
2. By two applications of the axiom of assignment and the sequential execution rule
 - | $\{ (M+1)*F = (M+1)! \wedge 0 \leq M+1 \leq N \} M := M+1; F := M*F \mathbb{P}$
3. Since $\mathbb{P} \wedge M < N \Rightarrow (M+1)*F = (M+1)! \wedge 0 \leq M+1 \leq N$, by the rule for strengthening the pre-condition, \mathbb{P} is an invariant for the while loop.

4. By step 3 and the while rule

\mathbb{P}
while $M < N$ **do**
begin $M := M+1; F := M * F$ **end**
 $\mathbb{P} \wedge M \geq N$

5. Since $\mathbb{P} \wedge M \geq N \Rightarrow F = N!$, by the rule for weakening the post-condition, the program proof is complete.