

## Equality and Rewriting Variations

We have already explored CafeOBJ's use of the rewriting relation  $\Rightarrow^*$  in place of equality, and some of the considerations that arise. CafeOBJ provides three other alternatives for term relationships (plus another 'is' for sorts). These relations are:

- $\_==\_$  for equality.
- $\_==>\_$  for transitions, and
- $\_=*=\_$ ,  $\_b=\_$  for behavioral equivalence (i.e., indistinguishability).

### Equality Predicate

The equality predicate ( $==$ ) is a vital operation of the system. It can be reliably used to test the equality of terms of the visible sorts when the  $\Rightarrow$  relation is confluent and terminating.

### Transition Predicate

The transition predicate ( $==>$ ) is an "oriented" version of equality for visible sorts. A transition relation is reflexive and transitive, but the symmetric property of equality ( $X=Y$  implies  $Y=X$ ) is omitted. Transitions are regarded as (normally) irreversible changes. However, the relation  $==>$  can be regarded as being defined by means of the following scheme of equality rules:

for each visible sort S             $\text{eq } X:S ==> X = \text{true}$   
for trans  $T \Rightarrow T1$              $\text{eq } T ==> T1 = \text{true}$   
for ctrans  $T \Rightarrow T1$  if C         $\text{ceq } T ==> T1 = \text{true if } C$

and for each

$\text{op } f : S1 \dots Sn \rightarrow S$          $\text{ceq } f(X1:S1, \dots, Xn:Sn) ==> f(Y1:S1, \dots, Yn:Sn) = \text{true}$   
   if  $X1 ==> Y1$  and ... and  $Xn ==> Yn$ .

This omits the transitive property. The direct way to express transitivity would be  $\text{ceq } X:S ==> X1:S = \text{true if } X ==> Y:S \text{ and } Y ==> X1$ .

However, such a rule involves a variable in the condition that does not appear in the left-hand side of the rule and is therefore prohibited. So instead CafeOBJ defines the operator " $\_=(*)=>\_$ " that is defined to mean a transition in an arbitrary number of steps, and the transitive property becomes the rule

$\text{ceq } X:S ==> Y:S = \text{true if } X =(*)=> Y$ .

### Behavioral Equivalence Predicate

Lastly, behavioral equivalence ( $=*==$ ) is defined for each hidden sort. Recall that behavioral operators (declared with **bop**) have exactly one argument with a hidden sort. The implication that two values of hidden sort H are indistinguishable is not fully captured by CafeOBJ's behavioral equivalence. The system uses only selectors (attributes as CafeOBJ calls them) with a single argument,  $f : H \rightarrow P$ , and if these are  $f1, f2, \dots, fn$ , then

$\text{eq } X:H =*== Y:H = f1(X) == f1(Y) \text{ and } \dots \text{ and } fn(X) == fn(Y)$ .

### **Evaluation alternatives**

There are three evaluation commands in CafeOBJ. They differ in the rules they employ during reduction, and in where they use them. In particular

- reduce and breduce use only equations, transitions are excluded,
- execute uses all the rules, and
- reduce and execute use behavioral rules only on subexpressions of a selector operation, while breduce uses them on all subexpressions.